# Total Colorings of Planar Graphs with Small Maximum Degree 

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#### Abstract

Let $G$ be a planar graph of maximum degree $\Delta$ and girth $g$, and there is an integer $t(>g)$ such that $G$ has no cycles of length from $g+1$ to $t$. Then the total chromatic number of $G$ is $\Delta+1$ if $(\Delta, g, t) \in\{(5,4,6),(4,4,17)\}$; or $\Delta=3$ and $(g, t) \in\{(5,13),(6,11),(7,11)$, $(8,10),(9,10)\}$, where each vertex is incident with at most one $g$-cycle.


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## 1. Introduction

All graphs considered in this paper are simple, finite and undirected, and we follow [2] for the terminologies and notations not defined here. Let $G$ be a graph. We use $V(G)$, $E(G), \Delta(G)$ and $\delta(G)$ (or simply $V, E, \Delta$ and $\delta$ ) to denote the vertex set, the edge set, the maximum degree and the minimum degree of $G$, respectively. For a vertex $v \in V$, let $N(v)$ denote the set of vertices adjacent to $v$, and let $d(v)=|N(v)|$ denote the degree of $v$. A $k$ vertex, a $k^{+}$-vertex or a $k^{-}$-vertex is a vertex of degree $k$, at least $k$ or at most $k$ respectively. A $k$-cycle is a cycle of length $k$, and a 3-cycle is usually called a triangle.

A total-k-coloring of a graph $G$ is a coloring of $V \cup E$ using $k$ colors such that no two adjacent or incident elements receive the same color. The total chromatic number $\chi^{\prime \prime}(G)$ of $G$ is the smallest integer $k$ such that $G$ has a total- $k$-coloring. Clearly, $\chi^{\prime \prime}(G) \geq \Delta+1$. Behzad [1] posed independently the following famous conjecture, which is known as the Total Coloring Conjecture (TCC).

Conjecture 1.1. For any graph $G, \Delta+1 \leq \chi^{\prime \prime}(G) \leq \Delta+2$.
This conjecture was confirmed for a general graph with $\Delta \leq 5$. But for planar graph, the only open case is $\Delta=6$ (see $[11,15]$ ). Interestingly, planar graphs with high maximum degree allow a stronger assertion,that is, every planar graph with high maximum degree $\Delta$ is $(\Delta+1)$-totally-colorable. This result was first established in [3] for $\Delta \geq 14$, which was extended to $\Delta \geq 9$ (see [12]). For $4 \leq \Delta \leq 8$, it is not known whether that the assertion still

[^0]holds true. But there are many related results by adding girth restrictions, see $[7-10,13,14$, 16]. We present our new results in this paper.

Theorem 1.1. Let $G$ be a planar graph of maximum degree $\Delta$ and girth $g$, and there is an integer $t(>g)$ such that $G$ has no cycles of length from $g+1$ to $t$. Then the total chromatic number of $G$ is $\Delta+1$ if $(a)(\Delta, g, t)=(5,4,6)$ or $(b)(\Delta, g, t)=(4,4,17)$.

Borodin et al. [6] obtained that if a planar graph $G$ of maximum degree three contains no cycles of length from 3 to 9 , then $\chi^{\prime \prime}(G)=\Delta+1$. In the following, we make further efforts on the total-colorability of planar graph on the condition that $G$ contains some $k$-cycle, where $k \in(5, \cdots, 9)$. We get the following result.

Theorem 1.2. Let $G$ be a planar graph of maximum degree 3 and girth $g$, each vertex is incident with at most one $g$-cycle and there is an integer $t(>g)$ such that $G$ has no cycles of length from $g+1$ to $t$. Then $\chi^{\prime \prime}(G)=\Delta+1$ if one of the following conditions holds.
(a) $g=5$ and $t \geq 13$,
(b) $g=6$ and $t \geq 11$,
(c) $g=7$ and $t \geq 11$,
(d) $g=8$ and $t \geq 10$,
(e) $g=9$ and $t \geq 10$.

We will introduce some more notations and definitions here for convenience. Let $G=$ $(V, E, F)$ be a planar graph, where $F$ is the face set of $G$. The degree of a face $f$, denoted by $d(f)$, is the number of edges incident with it, where each cut-edge is counted twice. A $k$-face or a $k^{+}$-face is a face of degree $k$ or at least $k$, respectively. Let $n_{k}(v)$ be the number of $k$-vertices adjacent to $v$ and $n_{k}(f)$ be the number of $k$-vertices incident with $f$.

## 2. Proof of Theorem 1.1

Let $G$ be a minimal counterexample to Theorem 1.1 in terms of the number of vertices and edges. Then every proper subgraph of $G$ is $(\Delta+1)$-totally-colorable. Firstly, we investigate some structural properties of $G$,which will be used to derive the desired contradiction completing our proof.

Lemma 2.1. $G$ is 2-connected and hence, it has no vertices of degree 1 and the boundary $b(f)$ of each face $f$ in $G$ is exactly a cycle (i.e. $b(f)$ cannot pass through a vertex $v$ more than once).

Lemma 2.2. [5] G contains no edge $u v$ with $\min \{d(u), d(v)\} \leq\lfloor\Delta / 2\rfloor$ and $d(u)+d(v) \leq$ $\Delta+1$.

Lemma 2.3. [3] The subgraph of $G$ induced by all edges joining 2-vertices to $\Delta$-vertices is a forest.

Lemma 2.4. [6] If $\Delta \geq 5$, then no 3-vertex is adjacent two 3-vertices.
Let $G_{2}$ be the subgraph induced by all edges incident with 2-vertices of $G$. Then $G_{2}$ is a forest by Lemma 2.3. We root it at a 5 -vertex. In this case, every 2 -vertex has exactly one parent and exactly one child, which are 5 -vertices.

Since $G$ is a planar graph, by Euler's formula, we have

$$
\sum_{v \in V}(d(v)-4)+\sum_{f \in F}(d(f)-4)=-8<0 .
$$

Now we define the initial charge function $\operatorname{ch}(x)$ of $x \in V \cup F$ to be $\operatorname{ch}(v)=d(v)-4$ if $v \in V$ and $\operatorname{ch}(f)=d(f)-4$ if $f \in F$. It follows that $\sum_{x \in V \cup F} \operatorname{ch}(x)<0$. Note that any
discharging procedure preserves the total charge of $G$. If we can define suitable discharging rules to change the initial charge function $c h$ to the final charge function $c h^{\prime}$ on $V \cup F$, such that $\operatorname{ch}^{\prime}(x) \geq 0$ for all $x \in V \cup F$, then we get an obvious contradiction. Now we design appropriate discharging rules and redistribute weights accordingly.

For a vertex $v$, we define $f_{k}(v)$ or $f_{k}^{+}(v)$ to be the number of $k$-faces or $k^{+}$-faces incident with $v$, respectively. To prove (a), our discharging rules are defined as follows.

R11. Each 2-vertex receives 2 from its child.
R12. Each 3-vertex $v$ receives $1 /\left(f_{7}^{+}(v)\right)$ from each of its incident $7^{+}$-faces.
R13. Each $5^{+}$-vertex receives $1 / 3$ from each of its incident $7^{+}$-faces.
Next, we will check $\operatorname{ch}^{\prime}(x) \geq 0$ for all $x \in V \cup F$. Let $f \in F(G)$. If $d(f)=4$, then $c h^{\prime}(f)=\operatorname{ch}(f)=0$. Suppose $d(f)=7$. Then $n_{3}(f) \leq 4$ by Lemma 2.4. Moreover, every $7^{+}$face sends at most $1 / 2$ to its incident 3 -vertices by R12 and $1 / 3$ to its incident 5 -vertices by R13. So we have $c h^{\prime}(f) \geq \operatorname{ch}(f)-4 \times 1 / 2-3 \times 1 / 3=0$. Suppose $d(f) \geq 8$. Then $n_{3}(v) \leq$ $\lfloor(2 d(f)) / 3\rfloor$ by Lemma 2.2 and Lemma 2.4. Thus, $c h^{\prime}(f) \geq c h(f)-(\lfloor(2 d(f)) / 3\rfloor \times 1 / 2)-$ $(d(f)-\lfloor(2 d(f)) / 3\rfloor) \times 1 / 3 \geq(5 d(f)-38) / 9 \geq 0$ by R12 and R13.

Let $v \in V(G)$. If $d(v)=2$, then $c h^{\prime}(v)=c h(v)+2=0$. If $d(v)=3$, then $f_{7}^{+}(v) \geq 2$ and it follows from R12 that $\operatorname{ch}^{\prime}(v)=\operatorname{ch}(v)+f_{7}^{+}(v) \times 1 /\left(f_{7}^{+}(v)\right)=0$. If $d(v)=4$, then $c h^{\prime}(v)=\operatorname{ch}(v)=0$. If $d(v)=5$, then $f_{7}^{+}(v) \geq 3$. Moreover, it may be the child of at most one 2 -vertex. Thus $c h^{\prime}(v) \geq \operatorname{ch}(v)+1 / 3 \times 3-2=0$ by R13. Suppose $d(v) \geq 6$. Then $v$ is incident with at most $\lfloor(d(v)) / 2\rfloor$ 4-faces and it may be the the parent of at most one 2-vertex. So $c h^{\prime}(v) \geq \operatorname{ch}(v)+(d(v)-\lfloor(d(v)) / 2\rfloor) \times 1 / 3-2=(7 d(v)-36) / 6>0$.

Note that (a) implies that (b) is true if $\Delta \geq 5$. So it suffice to prove (b) by assuming $\Delta=4$. Since $G$ is a planar graph, by Euler's formula, we have

$$
\sum_{v \in V}(2 d(v)-6)+\sum_{f \in F}(d(f)-6)=-12<0 .
$$

Now we define the initial charge function $\operatorname{ch}(x)$ of $x \in V \cup F$ to be $\operatorname{ch}(v)=2 d(v)-6$ if $v \in V$ and $\operatorname{ch}(f)=d(f)-6$ if $f \in F$. It follows that $\sum_{x \in V \cup F} c h(x)<0$.

To prove (b), we construct the new charge $c h^{\prime}(x)$ on $G$ as follows.
R21. Each $d(f)(d(f) \geq 18)$-face gives $1-6 /(d(f))$ to its incident vertices.
R22. Each 2 -vertex gets $3 / 2$ from its child and $1 / 2$ from its parent.
R23. Let $f$ be a 4 -face. If $f$ is incident with a 2 -vertex, then it gets $2 / 3$ from each of its incident $3^{+}$-vertices. If $f$ is incident with no 2 -vertices, then it gets $1 / 2$ from each of its incident vertices.
The rest of this paper is devoted to checking $c h^{\prime}(x) \geq 0$ for all $x \in V \cup F$. Let $f \in F(G)$. If $d(f)=4$, then $c h^{\prime}(f)=\operatorname{ch}(f)+\max \{2 / 3 \times 3,1 / 2 \times 4\}=0$. If $d(f) \geq 18$, then $c h^{\prime}(f)=$ $\operatorname{ch}(f)-r \times(1-6 / r)=0$ by R21.

Let $v \in V(G)$. If $d(v)=2$, then $c h^{\prime}(v)=\operatorname{ch}(v)+3 / 2+1 / 2=0$ by R22. If $d(v)=3$, then $f_{18}^{+}(v) \geq 2$ and $f_{4}(v) \leq 1$, and it follows from R21 and R23 that $c^{\prime}(v)=c h(v)+2 \times 2 / 3-$ $2 / 3>0$. Suppose that $d(v)=4$. Then $\operatorname{ch}(v)=2 \times 4-6=2$. If $n_{2}(v) \geq 1$, then $v$ sends at most $\left(n_{2}(v)+2\right) / 2$ to all its adjacent 2 -vertices by R22. If $3 \leq n_{2}(v) \leq 4$, then $f_{4}(v) \leq 1$ by Lemma 2.3, and it follows that $c h^{\prime}(v) \geq c h(v)-\left(n_{2}(v)+2\right) / 2+2 / 3 \times 3-2 / 3=(14-$ $\left.n_{2}(v) \times 3\right) / 6>0$ by R21 and R23. If $1 \leq n_{2}(v) \leq 2$, then $f_{4}(v) \leq 2$, and it follows that $\operatorname{ch}^{\prime}(v) \geq \operatorname{ch}(v)-\left(n_{2}(v)+2\right) / 2+2 / 3 \times 2-2 / 3 \times 2=\left(2-n_{2}(v)\right) / 2 \geq 0$. If $n_{2}(v)=0$, we have $f_{4}(v) \leq 2$. Moreover, each 4 -face incident with $v$ contains no 2 -vertices. By R23, we have $c^{\prime}(v) \geq \operatorname{ch}(v)+2 / 3 \times 2-1 / 2 \times 2>0$. Now we complete the proof of Theorem 1.1.

## 3. Proof of Theorem 1.2

A $3(k)$-vertex is a 3-vertex adjacent to exactly $k 2$-vertices. Let $G$ be a minimal counterexample to Theorem 1.2 in terms of the number of vertices and edges. By minimality of $G$, it has the following result.

Lemma 3.1. [6]
(a) no 2-vertex is adjacent to two 2-vertices;
(b) no 2-vertex is adjacent to a 2-vertex and a 3(2)-vertex;
(c) no 3-vertex is adjacent to three 2 -vertices.

Let $G_{23}$ be the bipartite subgraph of $G$ comprising $V$ and all edges of $G$ that join a 2-vertex to a 3-vertex. Then $G_{23}$ has no isolated 2-vertices by Lemma 3.1(a), and the maximum degree is at most 2 by Lemma 3.1(c), and any component of $G_{23}$ is a path with more than one edges must end in two 3 -vertices by Lemma 3.1(b). It follows that $n_{3} \geq n_{2}$. So we can find a matching $M$ in $G$ saturating all 2-vertices. If $u v \in M$ and $d(u)=2, v$ is called the 2 -master of $u$. Each 2 -vertex has one 2-master and each vertex of degree $\Delta$ can be the 2 -master of at most one 2-vertex.

Since $G$ is a planar graph, by Euler's formula, we have

$$
\sum_{v \in V}(d(v)-6)+\sum_{f \in F}(2 d(f)-6)=-12<0 .
$$

Now we define the initial charge function $\operatorname{ch}(x)$ of $x \in V \cup F$ to be $\operatorname{ch}(v)=d(v)-6$ if $v \in V$ and $\operatorname{ch}(f)=2 d(f)-6$ if $f \in F$. It follows that $\sum_{x \in V \cup F} \operatorname{ch}(x)<0$. Note that any discharging procedure preserves the total charge of $G$. If we can define suitable discharging rules to change the initial charge function $c h$ to the final charge function $c h^{\prime}$ on $V \cup F$, such that $c h^{\prime}(x) \geq 0$ for all $x \in V \cup F$, then we get an obvious contradiction. Now we design appropriate discharging rules and redistribute weights accordingly.
R31. Each $d(f)(d(f) \geq 5)$-face gives $2-6 /(d(f))$ to its incident vertices.
R32. Each 2-vertex receives $3-12 /(t+1)-6 / g$ from its 2-master.
Let $c h^{\prime}(x)$ be the new charge obtained by the above rules for all $x \in V \cup F$. If $f \in F(G)$, then $c^{\prime}(f)=\operatorname{ch}(f)-d(f) \times(2 d(f)-6) /(d(f))=0$ by R31. Let $v \in V(G)$. Suppose $d(v)=3$. Then $v$ can be the 2-master of at most one 2-vertex, and $v$ sends at most $3-$ $12 /(t+1)-6 / g$ to 2 -vertex by R32. In addition, If $v$ is incident with a $g$-face, then the other faces incident with $v$ are two $(t+1)^{+}$-faces, for $G$ has no cycles of length from $g+1$ to $t$. Thus, $v$ receives $(2-6 / g)$ from its incident $g$-face and $(2-6 /(t+1))$ from each of its incident $(t+1)^{+}$-face by R31. So $c h^{\prime}(v) \geq c h(v)+2(2-6 /(t+1))+(2-6 / g)-(3-$ $12 /(t+1)-6 / g)=0$ for all $g$ and $t$. Otherwise, $v$ is incident with three $(t+1)^{+}$-faces, then $c h^{\prime}(v) \geq \operatorname{ch}(v)+3(2-6 /(t+1))-(3-12 /(t+1)-6 / g)=6 / g-6 /(t+1)>0$, for $t+1>g$. Suppose $d(v)=2$. Then $v$ receives at most $3-12 /(t+1)-6 / g$ from its 2-master by R31. If $v$ is incident with a $g$-face, since $G$ has no cycles of length from $g+1$ to $t$, then the other face incident with $v$ is a $(t+1)^{+}$-face, and it follows that $c h^{\prime}(v) \geq \operatorname{ch}(v)+(2-$ $6 /(t+1))+(2-6 / g)+(3-12 /(t+1)-6 / g)=0$ for all $g$ and $t$. Otherwise, $v$ is incident with two $(t+1)^{+}$-faces, then $c h^{\prime}(v) \geq \operatorname{ch}(v)+2(2-6 /(t+1))+(3-12 /(t+1)-6 / g)=$ $3-24 /(t+1)-6 / g>0$.

From the above, we can see that $c^{\prime}(f)=c h(f)-d(f) \times(2 d(f)-6) /(d(f))=0$ for all $f \in F(G)$. Suppose $d(v)=3$. So $c h^{\prime}(v) \geq \operatorname{ch}(v)+2(2-6 /(t+1))+(2-6 / g)-(3-$ $12 /(t+1)-6 / g)=0$ for all $g$ and $t$. When $v$ is incident with three $(t+1)^{+}$-faces, then
$c h^{\prime}(v) \geq \operatorname{ch}(v)+3(2-6 /(t+1))-(3-12 /(t+1)-6 / g)=6 / g-6 /(t+1)>0$, for $t+1>$ $g$. Suppose $d(v)=2$. If $v$ is incident with a $g$-face and a $(t+1)^{+}$-face, then $c h^{\prime}(v) \geq \operatorname{ch}(v)+$ $(2-6 /(t+1))+(2-6 / g)+(3-12 /(t+1)-6 / g)=0$ for all $g$ and $t$. When $v$ is incident with two $(t+1)^{+}$-faces, then $c h^{\prime}(v) \geq c h(v)+2(2-6 /(t+1))+(3-12 /(t+1)-6 / g)=$ $3-24 /(t+1)-6 / g$. So when $g=5$, then $t \geq 13$; when $g=6$, then $t \geq 11$; when $g=7$, then $t \geq 11$; when $g=8$, then $t \geq 10$; when $g=9$, then $t \geq 10$, and it follows that $c h^{\prime}(v) \geq 0$.

Our proof of Theorem 1.2 is now complete.
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